Coarse Revealed Preference

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Revealed preference theory

Pioneered by Samuelson (1938)

Consumer theory: Afriat (1967)

General equilibrium theory: Brown and Matzkin (1996)

Industrial organization: Carvajal et al. (2013)

Matching theory: Echenique et al. (2013)

among many others.

Research question in its simplest form

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Investigate the observable restriction of economic models including

Rational choice with imperfect observation

Multiple preferences

Monotone multiple preferences

Minimax regret

Outline

Model

Theory

Applications

Related literature

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 B_i is a nonempty subset of A_i for each i.

To simplify the statements below, we write

$$C_i := A_i \setminus B_i$$
,

$$A(\mathcal{O}') := \cup_{(A_i,B_i)\in\mathcal{O}'} A_i$$

$$C(\mathcal{O}') := \cup_{(A_i,B_i)\in\mathcal{O}'} C_i$$
.

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Definition

A coarse data set $\mathcal O$ is coarsely rationalizable by a linear order if $\exists \ P$ such that

$$\max(A_i, P) \in B_i$$

for all i.

A coarse data set including four observations:

$$A_1 = \{x_1, x_2, x_3, x_4\}, B_1 = \{x_1, x_2\};$$

$$A_2 = \{x_2, x_3, x_4, x_5\}, B_2 = \{x_2, x_3\};$$

$$A_3 = \{x_3, x_4, x_5, x_1\}, B_3 = \{x_3, x_4\};$$

$$A_4 = \{x_4, x_5, x_1, x_2\}, B_4 = \{x_4, x_5\}.$$

Suppose that the data set is rationalizable by a linear order.

$$A_{1} = \{x_{1}, x_{2}, x_{3}, x_{4}\}, B_{1} = \{x_{1}, x_{2}\};$$

$$\implies (1a) x_{1}P^{*}x_{2}, x_{1}P^{*}x_{3}, x_{1}P^{*}x_{4}, \text{ or (1b) } x_{2}P^{*}x_{1}, x_{2}P^{*}x_{3}, x_{2}P^{*}x_{4};$$

$$A_{2} = \{x_{2}, x_{3}, x_{4}, x_{5}\}, B_{2} = \{x_{2}, x_{3}\};$$

$$A_{3} = \{x_{3}, x_{4}, x_{5}, x_{1}\}, B_{3} = \{x_{3}, x_{4}\};$$

$$A_{4} = \{x_{4}, x_{5}, x_{1}, x_{2}\}, B_{4} = \{x_{4}, x_{5}\}.$$

Suppose that the data set is coarsely rationalizable by a linear order.

(1a)
$$x_1P^*x_2$$
, $x_1P^*x_3$, $x_1P^*x_4$, or (1b) $x_2P^*x_1$, $x_2P^*x_3$, $x_2P^*x_4$;

(2a)
$$x_2P^*x_3$$
, $x_2P^*x_4$, $x_2P^*x_5$, or (2b) $x_3P^*x_2$, $x_3P^*x_4$, $x_3P^*x_5$;

(3a)
$$x_3P^*x_1$$
, $x_3P^*x_4$, $x_3P^*x_5$, or (3b) $x_4P^*x_1$, $x_4P^*x_3$, $x_4P^*x_5$;

(4a)
$$x_4P^*x_1$$
, $x_4P^*x_2$, $x_4P^*x_5$, or (4b) $x_5P^*x_1$, $x_5P^*x_2$, $x_5P^*x_4$.

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Considering all $2^{|\mathcal{O}|}$ possible combinations?

Suppose that $\ensuremath{\mathcal{O}}$ is coarsely rationalizable by a linear order.

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Consider any nonempty subcollection $\mathcal{O}' = \{(A_{k_j}, B_{k_j})\}_{j=1}^m$ of \mathcal{O} .

 \implies the maximal element in A_{k_j} is not contained in C_{k_j} , $\forall j$

 \implies the maximal element in $A(\mathcal{O}')$ is not contained in $C(\mathcal{O}')$.

Necessary condition for coarse rationalizability that we call Coarse SARP:

Coarse SARP. For any $\emptyset \neq \mathcal{O}' \subseteq \mathcal{O}$, $A(\mathcal{O}') \setminus C(\mathcal{O}') \neq \emptyset$.

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$$A(\mathcal{O}) \setminus C(\mathcal{O}) = \emptyset$$

⇒ Violation of Coarse SARP

⇒ Not coarsely rationalizable by a linear order.

Coarse SARP is also a sufficient condition.

Theorem

A coarse data set is coarsely rationalizable by a linear order if and only if it satisfies the Coarse SARP property.

Illustrating the proof using an example

Example

Consider the following coarse data set including five observations:

	A _i	B _i	C _i
i = 1	$\{x_1, x_2, x_3, x_4\}$	$\{x_1,x_2\}$	$\{x_3, x_4\}$
<i>i</i> = 2	$\{x_2, x_3, x_4, x_5\}$	$\{x_2, x_3\}$	$\{x_4, x_5\}$
<i>i</i> = 3	$\{x_3, x_4, x_5, x_6\}$	$\{x_3, x_4\}$	$\{x_5, x_6\}$
i = 4	$\{x_4, x_5, x_6, x_7\}$	$\{x_4, x_5\}$	$\{x_6, x_7\}$
<i>i</i> = 5	$\{x_5, x_6, x_7, x_1\}$	$\{x_5, x_6, x_7\}$	$\{x_1\}$

Let $\mathcal{O}_1 := \mathcal{O}$.

	Ai	B _i	Ci
i = 1	$\{x_1, x_2, x_3, x_4\}$	$\{x_1,x_2\}$	$\{x_3, x_4\}$
i=2	$\{x_2, x_3, x_4, x_5\}$	$\{x_2,x_3\}$	$\{x_4,x_5\}$
i = 3	$\{x_3, x_4, x_5, x_6\}$	$\{x_3, x_4\}$	$\{x_5, x_6\}$
i = 4	$\{x_4, x_5, x_6, x_7\}$	$\{x_4, x_5\}$	$\{x_6, x_7\}$
i = 5	$\{x_5, x_6, x_7, x_1\}$	$\{x_5, x_6, x_7\}$	{ <i>x</i> ₁ }

Then $A(\mathcal{O}_1) \setminus C(\mathcal{O}_1) = \{x_2\}.$

Let $P_1 := A(\mathcal{O}_1) \setminus C(\mathcal{O}_1) = \{x_2\}.$

Rank x above y if $x \in P_1$ and $y \in A(\mathcal{O}) \setminus P_1$.

Let $\mathcal{O}_2 := \{(A_i, B_i) \in \mathcal{O}_1 : A_i \cap P_1 = \emptyset\}.$

	A_i	B _i	Ci
i = 3	$\{x_3, x_4, x_5, x_6\}$	$\{x_3, x_4\}$	$\{x_5, x_6\}$
i = 4	$\{x_4, x_5, x_6, x_7\}$	$\{x_4, x_5\}$	$\{x_6, x_7\}$
<i>i</i> = 5	$\{x_5, x_6, x_7, x_1\}$	$\{x_5, x_6, x_7\}$	$\{x_1\}$

Repeat this logic...

 \mathcal{O} is finite...

Strict partial order \rightarrow linear order.

Coarse SARP and the classical SARP

In the special case that B_i is a singleton set for each i,

Coarse SARP reduces to the classcial SARP.

Both directions are easy to verify.

Application 1: Rational choice with imperfect observation

We represent the observed behavior of the DM by (Σ, f) , where

$$\Sigma \subset \mathcal{X}$$
,

f(A) is superset of the choice of the DM in $A \in \Sigma$.

Application 2: Multiple preferences

The DM has a set ▷ of strict preferences, and she chooses

$$f_{\triangleright}(A) := \{x \in A : x = \max(A, \succ) \text{ for some } \succ \in \triangleright\}$$

from each feasible set A.

See, for example, Salant and Rubinstein (2008).

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f(A) is the set of all alternatives that the DM chooses in $A \in \Sigma$.

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We say that (Σ, f) is rationalizable by multiple preferences if there exists a set \triangleright of strict preferences such that

$$f_{\triangleright}(A) = f(A)$$

for all $A \in \Sigma$.

Divide and conquer

For each $A \in \Sigma$ and $x \in f(A)$, we construct a coarse data set $\mathcal{O}_{A,x}$ indexed by (A,x) as follows:

$$\mathcal{O}_{A,x}:=\{(A',f(A'))\}_{A'\in\Sigma,A'\neq A}\cup(A,x).$$

Let

$$\mathfrak{D} := \{\mathcal{O}_{A,x}\}_{A \in \Sigma, x \in f(A)}.$$

A necessary condition for the data set (Σ, f) to be rationalizable by multiple preferences is that each $\mathcal{O}_{A,x}$ constructed in this way is rationalizable by a linear order.

Theorem

 (Σ, f) is rationalizable by multiple preferences

if and only if

each $\mathcal{O}_{A,x}$ in $\mathfrak D$ is rationalizable by a linear order.

Application 3: Minimax regret

Let $u: X \to R$ be a utility function for the DM.

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Given a finite set of utility functions \mathcal{U} , the worst-case regret of choosing x from $A \in \mathcal{X}$ is

$$\max_{y \in A} \max_{u \in \mathcal{U}} \left[u(y) - u(x) \right].$$

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The DM has a finite set of utility functions $\mathcal U$ defined on X and she chooses

$$\min_{x \in A} \left\{ \max_{y \in A} \max_{u \in \mathcal{U}} \left[u(y) - u(x) \right] \right\}.$$

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We say that (Σ, f) is rationalizable under the minimax regret model if there is a finite set of utility functions \mathcal{U} such that

$$f(A) = \underset{x \in A}{\operatorname{arg \, min}} \left\{ \max_{y \in A} \max_{u \in \mathcal{U}} \left[u(y) - u(x) \right] \right\}.$$

for each $A \in \Sigma$.

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for each $A \in \Sigma$.

For simplicity, write $\phi(x, y) = \max_{u \in \mathcal{U}} [u(y) - u(x)]$.

Suppose that (Σ, f) includes the following observation $f(\{x, y, z\}) = x$.

It must be the case that

$$\max \{\phi(y,x),\phi(y,z)\} > \max \{\phi(x,y),\phi(x,z)\}$$

and

$$\max \{\phi(z,x),\phi(z,y)\} > \max \{\phi(x,y),\phi(x,z)\}.$$

Construct a corresponding coarse data set...

Related Literature

Fishburn (1976)

Partial congruence axiom

de Clippel and Rosen (2018)

Bounded rationality theories under incomplete data

Enumeration procedure

Hu et al. (2018)

Explore related ideas in different settings

Weak order